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# **Characterizing One-/Two-sided Designs in OpenSHMEM Collectives**

Yuke Li<sup>1</sup>, Yanfei Guo<sup>2</sup>, and Xiaoyi Lu<sup>1</sup>

<sup>1</sup>University of California Merced, {yli304, xiaoyi.lu}@ucmerced.edu <sup>2</sup>Argonne National Laboratory, yguo@anl.gov

# **Introduction and Motivation**

- OpenSHMEM is a widely used PGAS programming model in the HPC area
- The team concept and team-based collective communication in the latest **OpenSHMEM Specification v1.5 are** similar to the communicator and collective communication in MPI



# **Designs of OpenSHMEM Libraries**

- Sandia OpenSHMEM (SOS)
  - ✤ Native one-sided design -- a source PE puts data into (or gets data from) the shared global memory of the target PE
  - The synchronization of the one-sided design does not require the target PE to actively acknowledge data receiving to the source PE
  - **SOS collectives natively implement some** algorithms – Tree, Ring, Round-Robin algorithms



#### OpenSHMEM over MPI (OSHMPI)

MPI-based two-sided design -- a source PE sends data, which matches with a *recv* called by a target PE

Argonne 🕰

- Less flexible synchronization. It needs send and recv coordination between the source and the target PEs
- ✤ OSHMPI collectives essentially call MPI collectives. Taking MPICH as an example, it provides rich collective algorithms – Tree, Ring, Bruck, Recursive Doubling, etc. The algorithm can be chosen at runtime by scheduling



Shared memory	Shared memory	Shared memory	F	Global Address Space (PGAS)	
Memory space	Memory space	Memory space			

Fig. 1: An OpenSHMEM team

- The different design approaches (the native one-sided and the MPI-based two-sided communication) can lead to different performance characteristics on HPC clusters (even for the same collective routine)
- We characterize two aspects that can influence the performance
  - Synchronization methods
  - Collective algorithms
- We compare the native one-sided design and the MPI-based two-sided

data	target		Partitioned Global Address Space (PGAS)
Memory space	Memory space	Memory space	

Fig. 2: One-sided communication design -- A PE *puts* data directly to target shared global memory of the other PE

data	target		Partitioned Global Address Space (PGAS)
Memory space	Memory space	Memory space	

Fig. 3: Two-sided communication design -- A send called by source PE must match with a *recv* called by target PE

### **Performance Characterization**







design in collectives using benchmarks, and show how big the performance difference between them

# **Experiment Setup**

• Library Selection:

One-sided design

Sandia OpenSHMEM (SOS) [1]

Two-sided design

**OpenSHMEM over MPI (OSHMPI) [2][3]** 

• Platform:

- **\*** Bebop HPC cluster [4]. Each node has **36-core Intel Broadwell CPU. The nodes** are connected with Omni-Path Fabric
- MPI Implementation: MPICH [5][6]
- **Benchmark**: OSU Micro-Benchmarks [7]

Key Findings	Reason Explanations
OSHMPI (two-sided) is faster than SOS (one-sided) in most cases.	The existing two-sided design (e.g., OSHMPI) can inherit the advantages in well-optimized MPI

- The performance of OpenSHMEM varies a lot (up to 10X) with several factors
- The comparison exposes the performance characteristics of different collective communication designs in OpenSHMEM

• We believe the performance characterizations can give the community some insights for future research avenues SOS shows lower or comparable latency in certain cases: - Communication with a small number of PEs like 2, for *collect*, intra-node broadcast, alltoall, and reduce collectives. - The *alltoall* collective for medium messages in intra-node communication and medium-large messages in internode communication with more PEs, like 16 and 32.

SOS shows unstable performance in the case of intra-node communication with 32 PEs.

engine, while the one-sided design (e.g., SOS) still needs optimized implementations.

With a small number of PEs, basic point-to-point primitive performance and synchronization method dominate the performance. SOS's native one-sided design can achieve better overlapping with simpler synchronizations. Besides, SOS also has some specific optimizations on Omni-Path Fabric. For the *alltoall* collective, MPI internal routine needs further optimization or tuning for these particular settings.

SOS uses extra helper threads to progress the internal communication tasks and thus it is oversubscribed.

### References

Conclusion

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